# LAVA: Large-scale Automated Vulnerability Addition

Tim Leek, Patrick Hulin, Ryan Whelan (MIT/LL),

Brendan Dolan-Gavitt (NYU),

Fredrick Ulrich, Andrea Mambretti, Wil Robertson, and Engin Kirda (Northeastern)





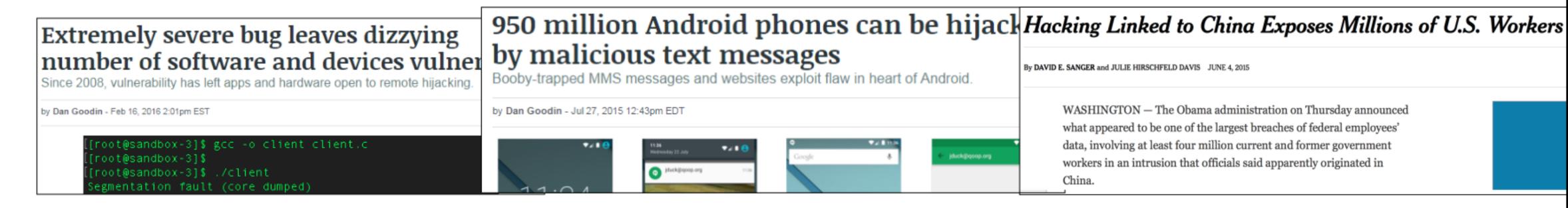








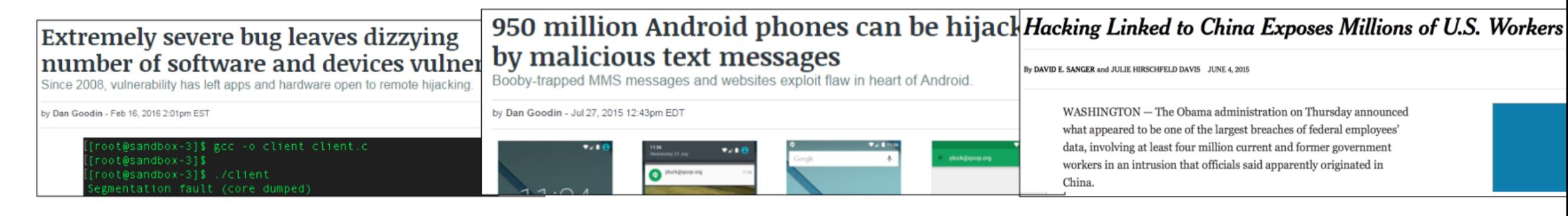
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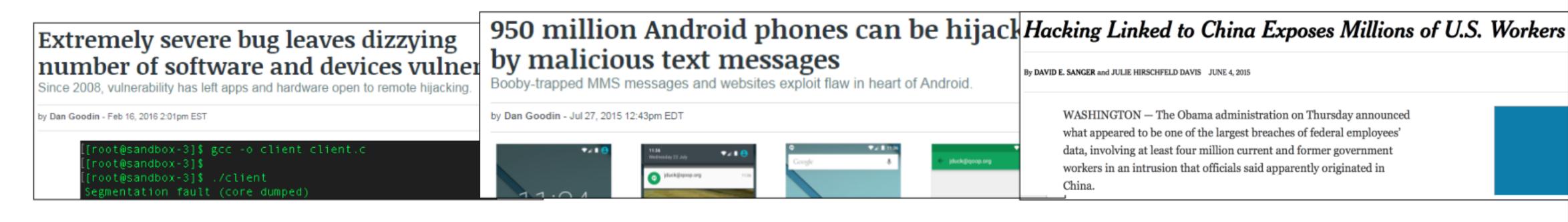


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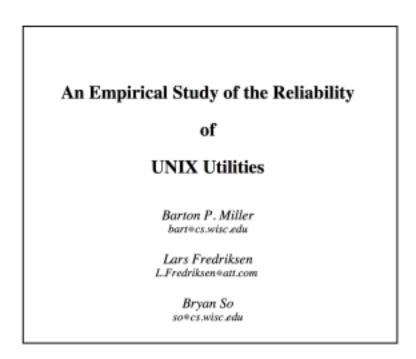




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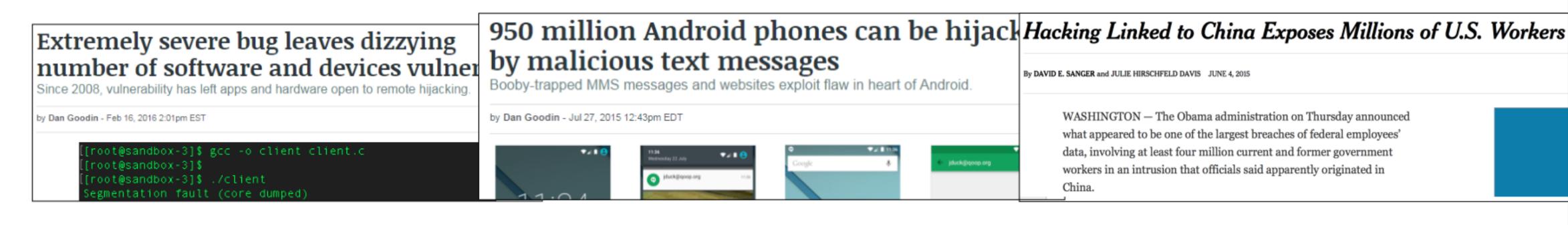


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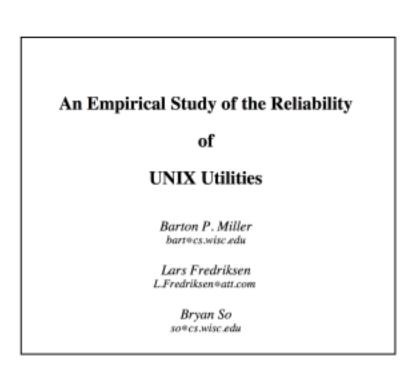




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A Functional Method for Assessing Protocol Implementation Security

Rauli Kaksonen

VTT Electronics

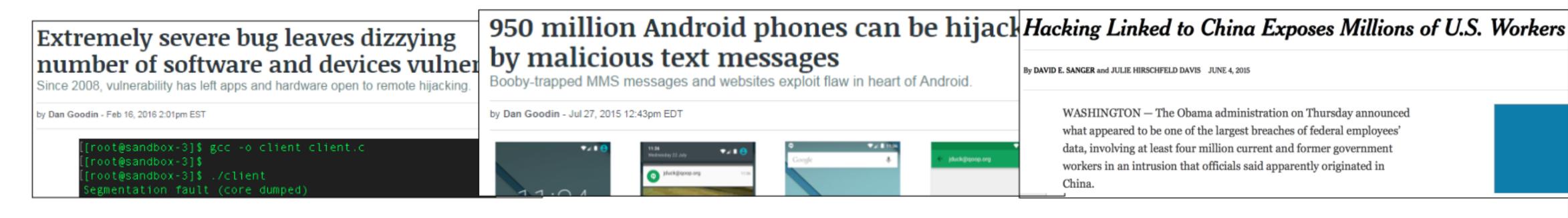
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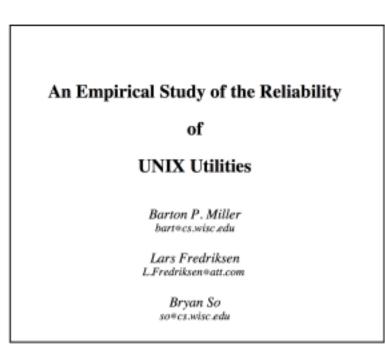




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Tests for Complex Systems Programs

Cristian Cadar, Daniel Dunbar, Dawson Engler \*
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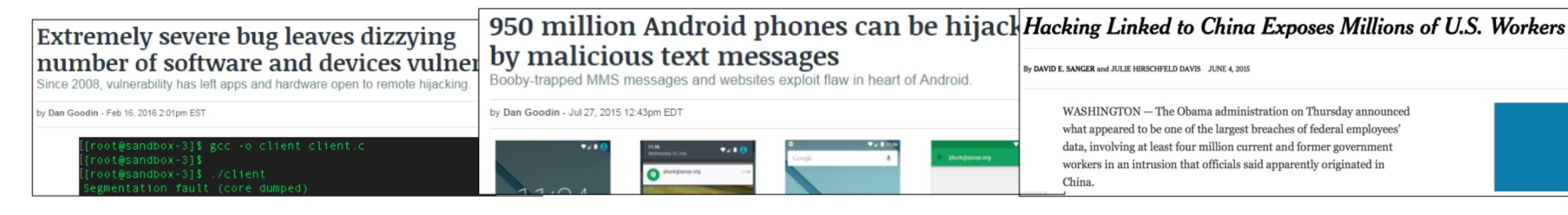
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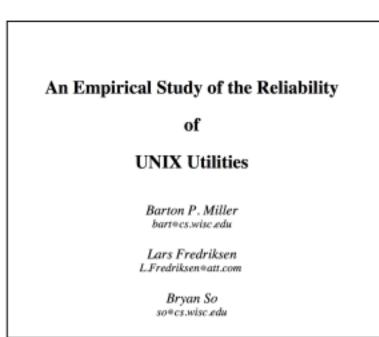




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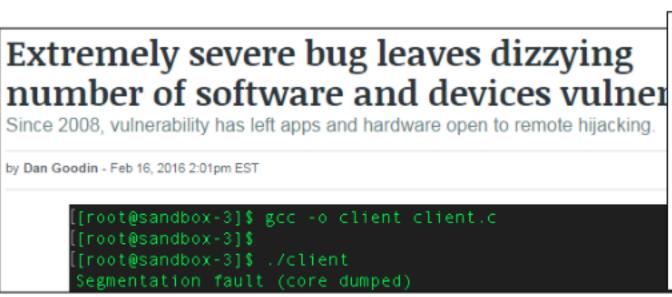
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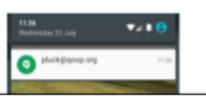


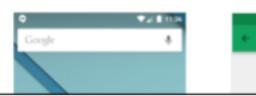
### 950 million Android phones can be hijacl Hacking Linked to China Exposes Millions of U.S. Workers by malicious text messages

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by Dan Goodin - Jul 27, 2015 12:43pm EDT



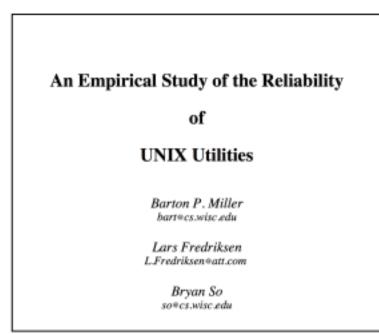




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INDUSTRY

Tim Leek- 2 TRL 02/25/16





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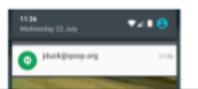
Extremely severe bug leaves dizzying number of software and devices vulner Since 2008, vulnerability has left apps and hardware open to remote hijacking. by Dan Goodin - Feb 16, 2016 2:01pm EST @sandbox-3]\$ gcc -o client client.c t@sandbox-3]\$ ./client mentation fault (core dumped)

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An Empirical Study of the Reliability **UNIX Utilities** Barton P. Miller hart@cs.wisc.edu Lars Fredriksen L.Fredriksen@att.com Bryan So so⊕cs.wisc.edu

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# Existing vulnerability corpora















# Existing vulnerability corpora













Testing Static Analysis Tools using Exploitable Buffer Overflows from Open Source Code

#### ABSTRAC

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Kendra Kratkiewicz MIT Lincoln Laboratory 244 Wood Street Lexington, MA 02420-9108 Phone: 781-981-2931 Email: KENDRA@LL.MIT.EDU

#### ABSTRACT

A corpus of 291 small C-program test cases was developed to evaluate static and dynamic analysis tools designed to detect buffer overflows. The corpus was designed and labeled using a new, comprehensive buffer overflow taxonomy. It provides a benchmark to measure detection, false alarm, and confusion rates of tools, and also suggests areas for tool enhancement. Experiments with five tools demonstrate that some modern static analysis tools can accurately detect overflows in simple test cases but that others have serious limitations. For example, PolySpace demonstrated a superior detection rate, missing only one detection. Its performance could be enhanced if extremely long run times were reduced, and false alarms were eliminated for some C library functions. ARCHER performed well with no false alarms whatsoever. It could be enhanced by improving interprocedural analysis and handling of C library functions. Splint

Richard Lippmann MIT Lincoln Laboratory 244 Wood Street Lexington, MA 02420-9108 Phone: 781-981-2711 Email: LIPPMANN@LL.MIT.EDU

for a significant percentage of the software vulnerabilities published each year [17, 19], such as in NIST's ICAT Metabase [8], CERT advisories [1], Bugtraq [16], and other security forums. Buffer overflows have also been the basis for many damaging exploits, such as the Sapphire/Slammer [12] and Blaster [14] worms.

A buffer overflow vulnerability occurs when data can be written outside the memory allocated for a buffer, either past the end or before the beginning. Buffer overflows may occur on the stack, on the heap, in the data segment, or the BSS segment (the memory area a program uses for uninitialized global data), and may overwrite from one to many bytes of memory outside the buffer. Even a one-byte overflow can be enough to allow an exploit [9]. Buffer overflows have been described at length in many papers, including [19], and many descriptions of exploiting buffer overflows can be found online.



# Existing vulnerability corpora



2005











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ADOBE READER	\$5,000-\$30,000
MAC OSX	\$20,000-\$50,000
ANDROID	\$30,000-\$60,000
FLASH OR JAVA BROWSER PLUG-INS	\$40,000-\$100,000
MICROSOFT WORD	\$50,000-\$100,000
WINDOWS	\$60,000-\$120,000
FIREFOX OR SAFARI	\$60,000-\$150,000
CHROME OR INTERNET EXPLORER	\$80,000-\$200,000
IOS	\$100,000-\$250,000

Forbes, 2012



# Vulnerability corpora sources



Source	Cost	Realism	Yield	
Accident	FREE	High	Tiny	
Search	\$\$\$\$	Med-High	Low	
Injection	\$\$	Med	Low-Med	
Synthesis	\$	Low	High	









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# LAVA concept



- Vulnerability corpus requirements
  - Cheap and plentiful
  - ☐ Realistic
  - ☐ Triggering input
  - ☐ Manifest only for one or very few inputs
  - ☐ Security-critical effect



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- Caveats
  - Works only on source
  - C programs
  - Linux
  - Buffer overflows



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- Large-scale Automated Vulnerability Addition
  - Uses static and dynamic analysis to find attacker-controlled data that can be used to introduce new code that creates a bug
  - Change program and input at same time to insert bugs in known places
  - Special sauce: new taint-based measures



# Dynamic taint analysis

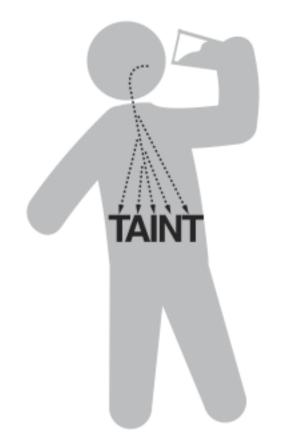


### PANDA dynamic taint

- Whole system (all processes + kernel)
- Works on binaries
- Includes all library code
- Oddball x86 instructions all analyzed including FPU and SSE
- Many labels supported: Every byte in 10MB file
- Labels combine into sets to represent computation
- Fast (enough). 50-100x



https://github.com/panda-re

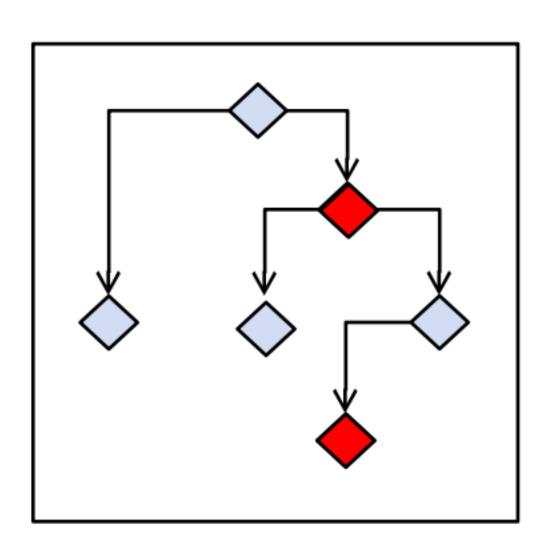










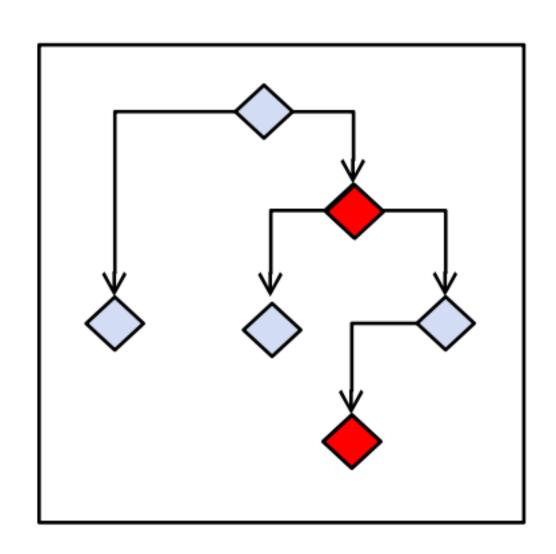


#### <u>Liveness</u>:

Number of branches an input byte is used to decide.
How much effect upon control flow do specific input bytes have?

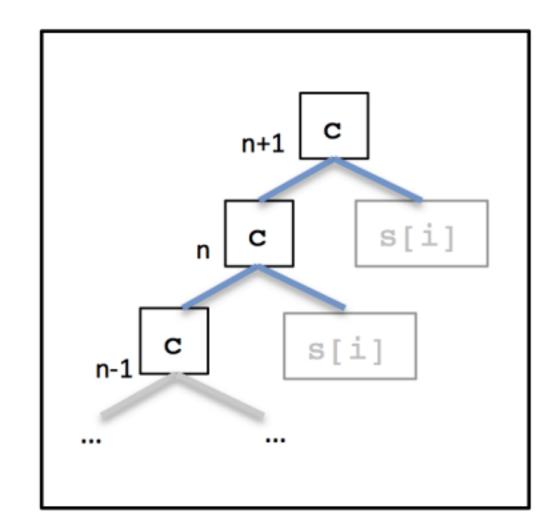






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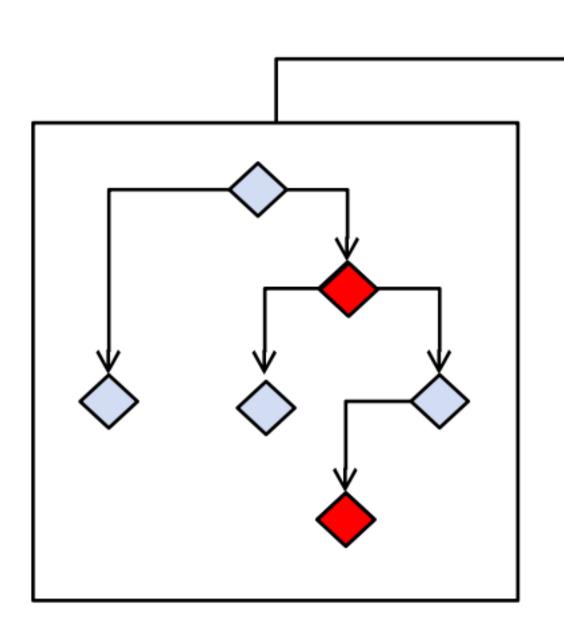


### Taint compute number:

Depth of Ival tree of computation. How complicated a function of input bytes is an Ival?







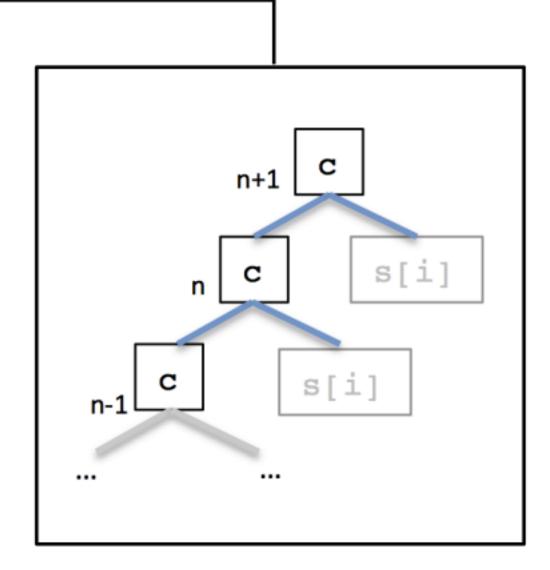
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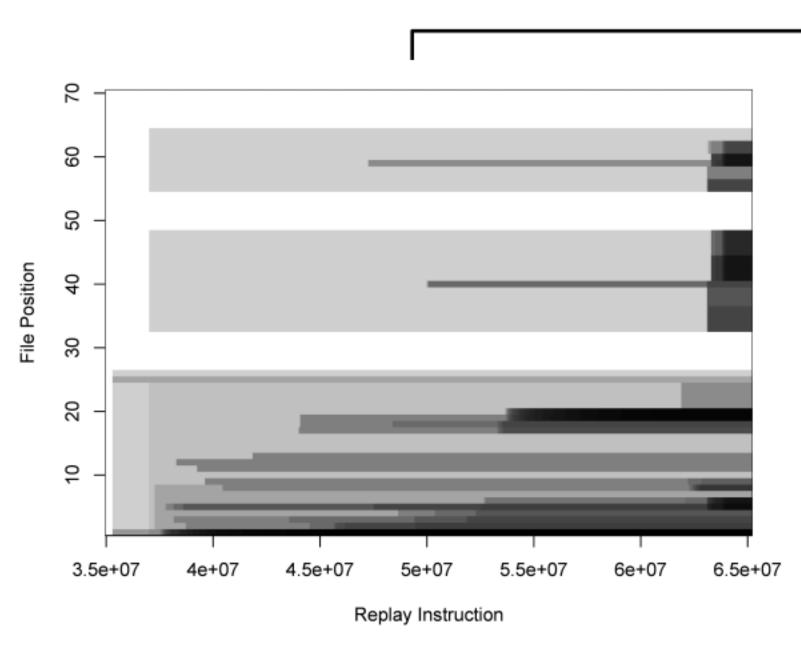
Attacker-controlled data that can be used to create a vulnerability



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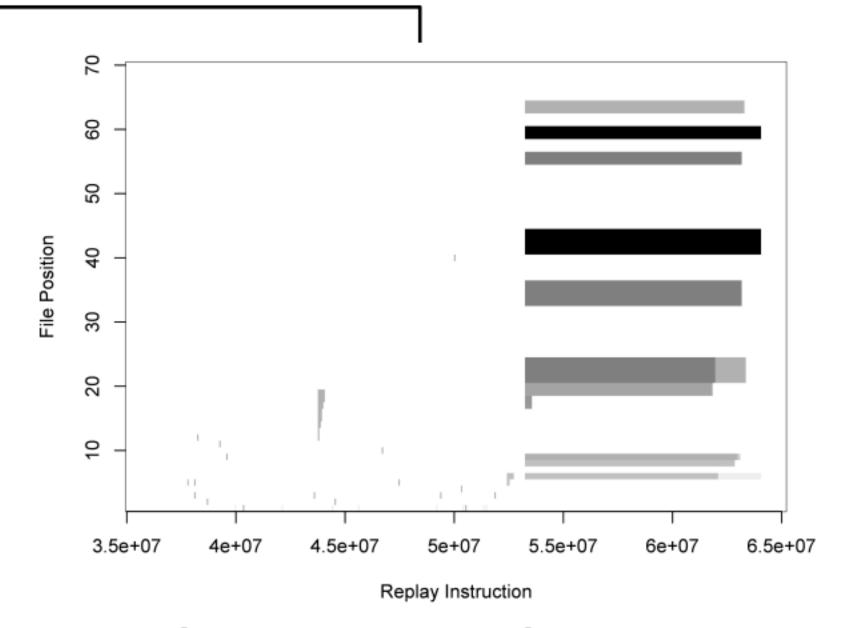
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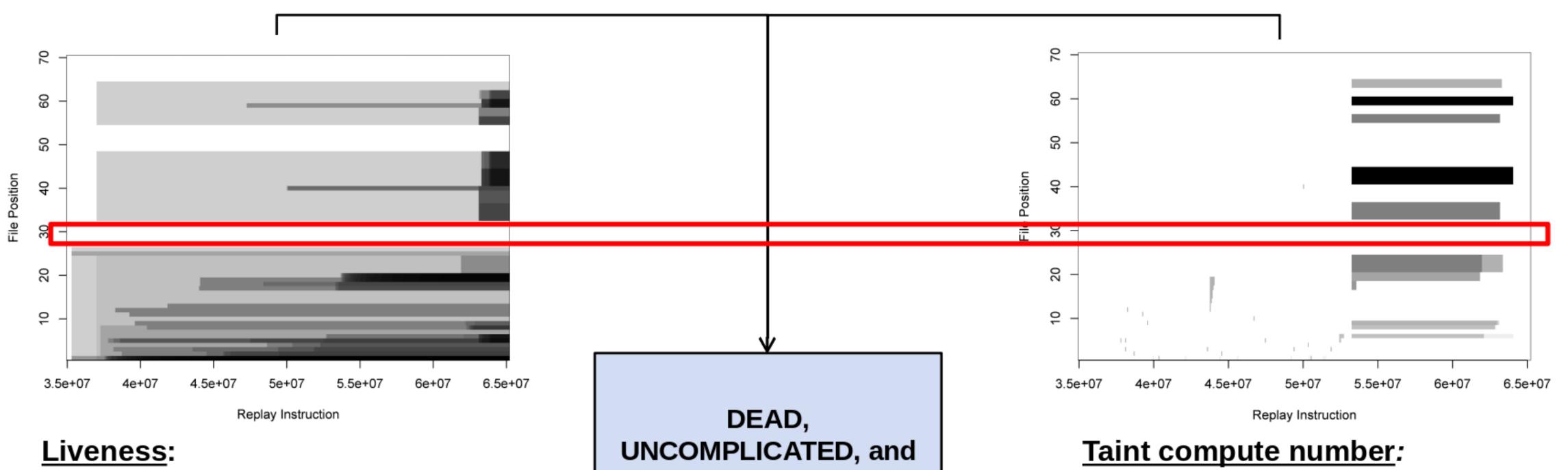
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**AVAILABLE data (DUA)** 

**Attacker-controlled data** that can be used to create a vulnerability

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Clang

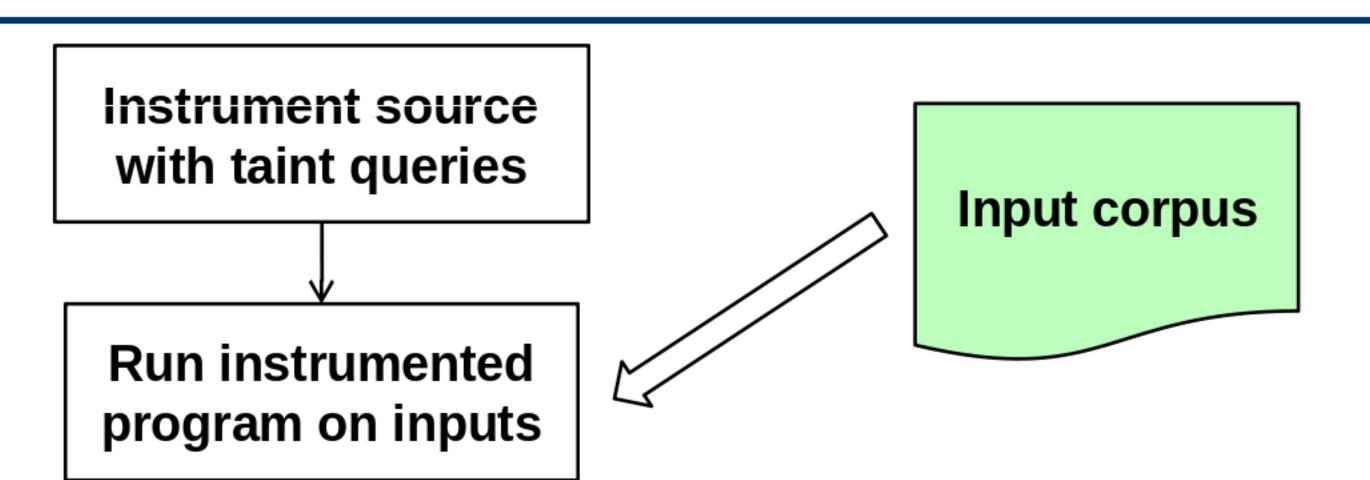
Instrument source with taint queries





Clang

**PANDA** record



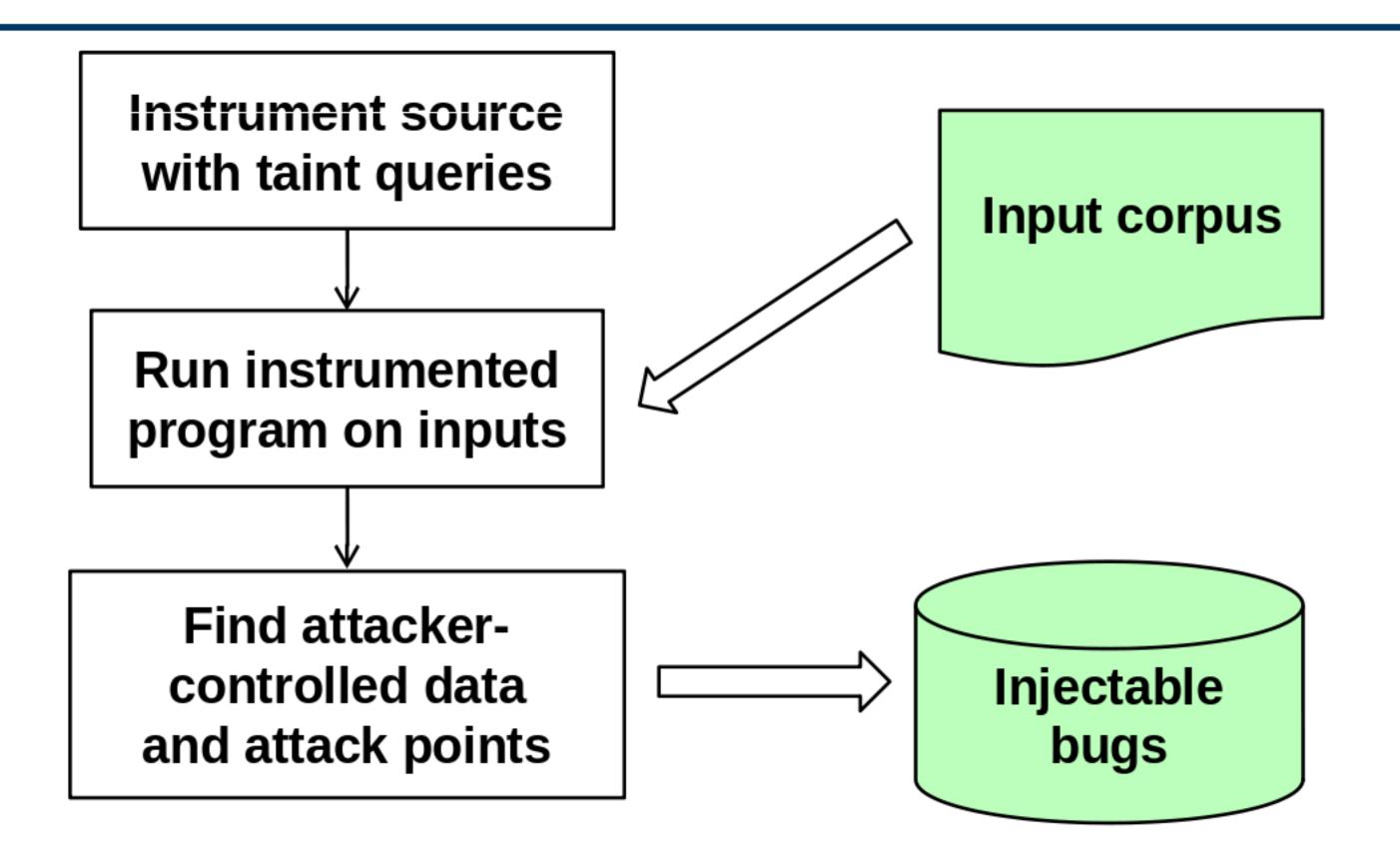




Clang

**PANDA** record

PANDA replay + taint analysis





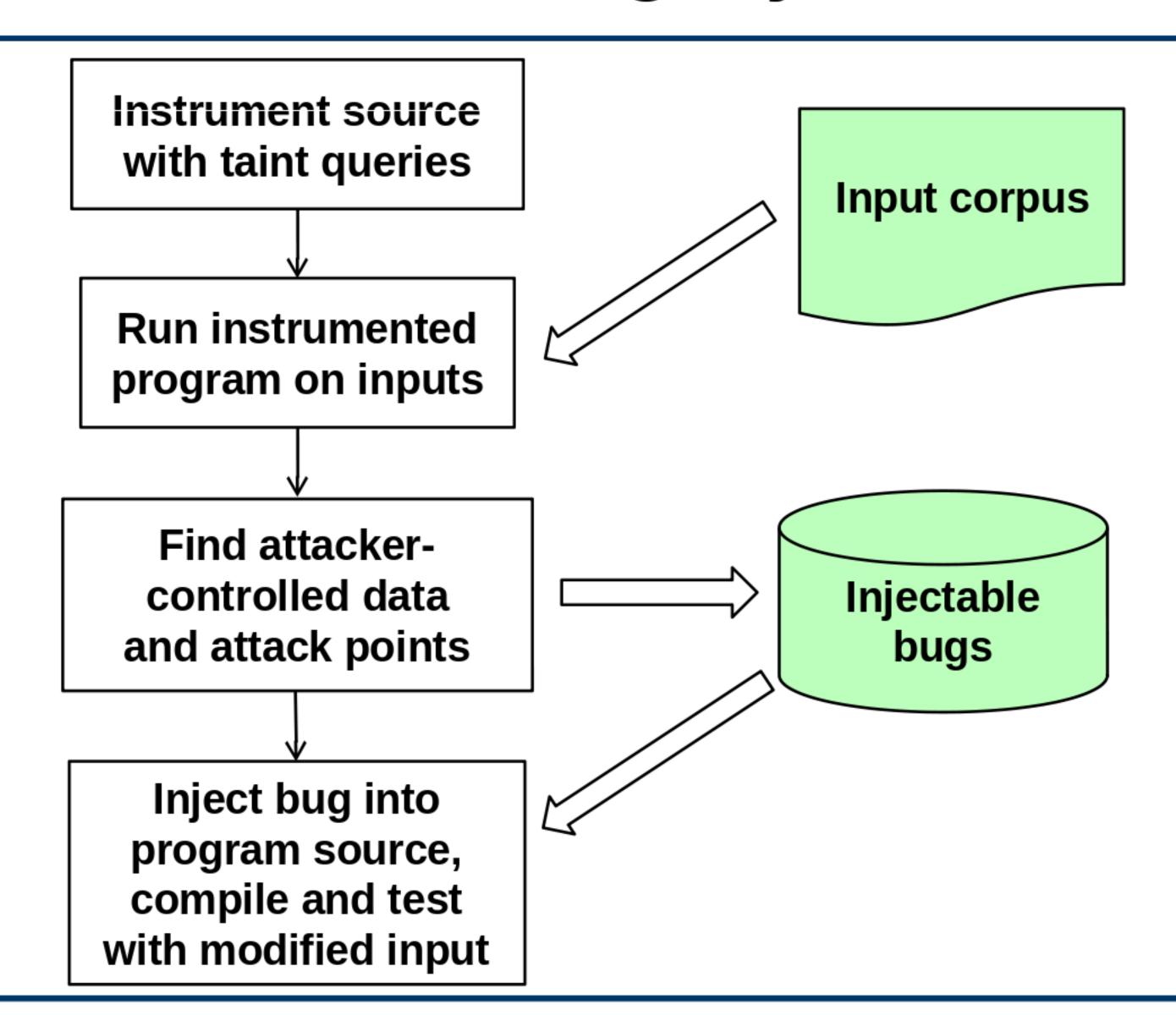


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**PANDA** record

PANDA replay + taint analysis

Clang





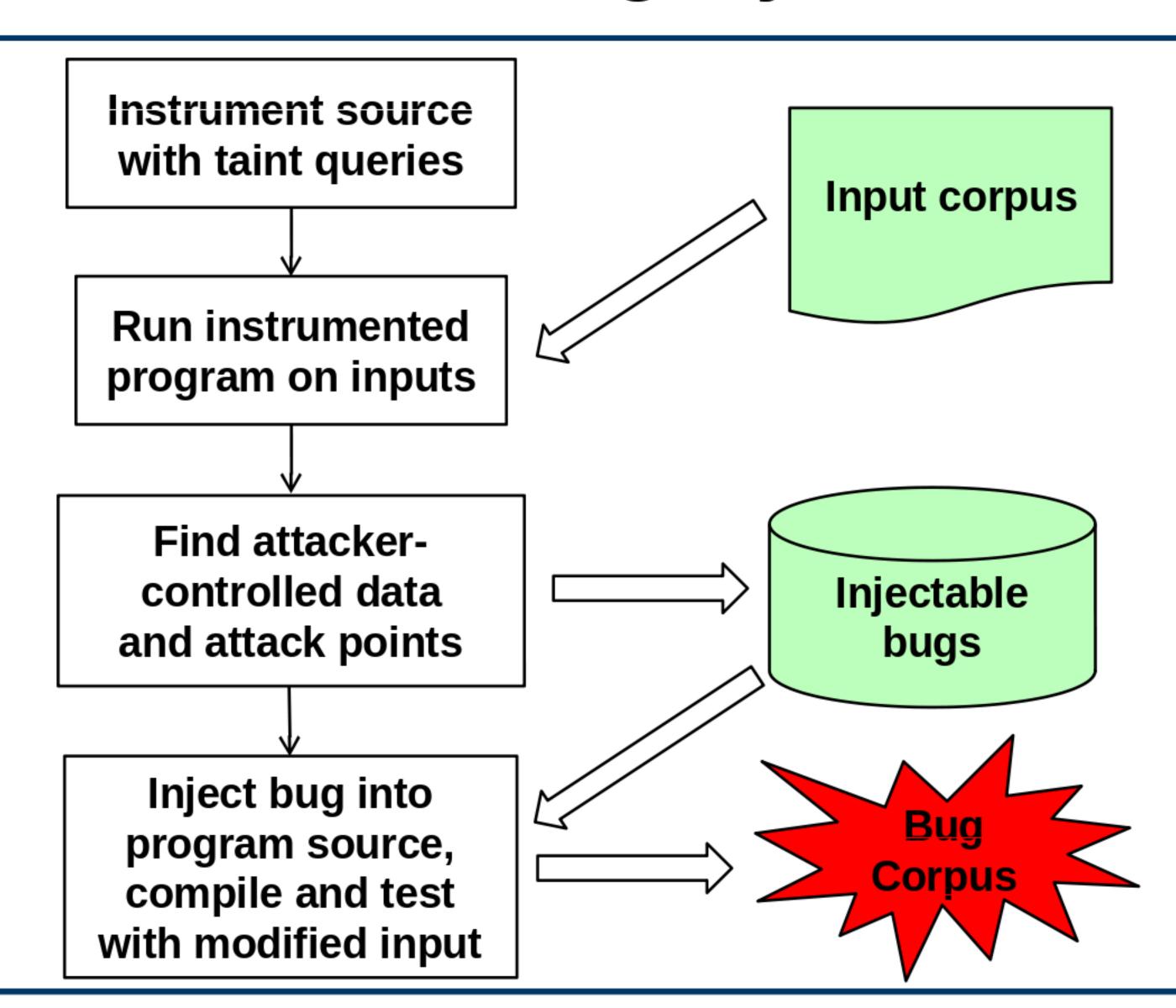


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 PANDA taint analysis tells us that bytes 0-3 in the buffer buf at line 115 of src/encoding.c is attacker-controlled





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### Attacker controlled data





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Attacker controlled data
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Corruptible pointer

readcdf.c 365: if (cdf\_read\_header(&info, &h) == -1)





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Corruptible pointer

**New data flow** 

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# Vulnerability injection effectiveness



TABLE I

LAVA INJECTION RESULTS FOR OPEN SOURCE PROGRAMS OF VARIOUS SIZES

		Num	Lines			Potential	Validated		Inj Time
Name	Version	Src Files	C code	N(DUA)	N(ATP)	Bugs	Bugs	Yield	(sec)
file	5.22	19	10809	631	114	17518	774	38.7%	16
readelf	2.25	12	21052	3849	266	276367	1064	53.2 %	354
bash	4.3	143	98871	3832	604	447645	192	9.6%	153
tshark	1.8.2	1272	2186252	9853	1037	1240777	354	17.7%	542

- Four open source programs 10K -> 2M LOC
- 2000 injection attempts per target (of over 1M)
- LAVA yield (validated injected bugs): 10->50%
- Over 2000 bugs injected



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Over 200K possible?

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# Using LAVA to evaluate tools



- Created two corpora using LAVA
  - LAVA-1 programs containing individual bugs of varying difficulty
  - LAVA-M programs each with more than one bug

- Evaluated two open-source vulnerability discovery tools by ability to detect LAVA bugs
  - Fuzzer
  - Symbolic execution + SAT solving

TABLE IV
BUGS FOUND IN *LAVA-M* CORPUS BY TOOL TYPE

Tool Name	Total Bugs	Unique Bugs Found					
1001 Name	Total Bugs	FUZZER	SES	Combined			
uniq	28	7	0	7			
base64	44	7	9	14			
md5sum	57	2	0	2			
who	2136	0	18	18			
Total	2265	16	27	41			



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Detection < 2%



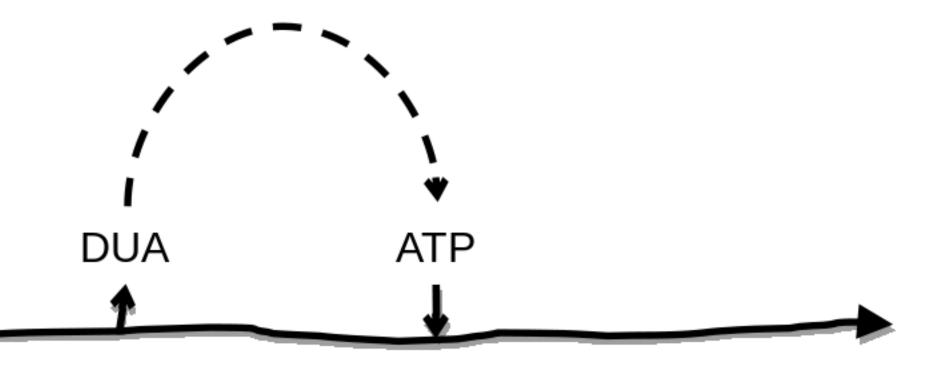
# LAVA vulnerability realism



Realism is a concern. But hard to quantify

One possible measure is the fraction of the trace that is unaffected by LAVA yet must be analyzed correctly to discover the vulnerability

LAVA's bugs are inserted, generally quite far along in the trace. If anything we need some easier ones



Execution trace

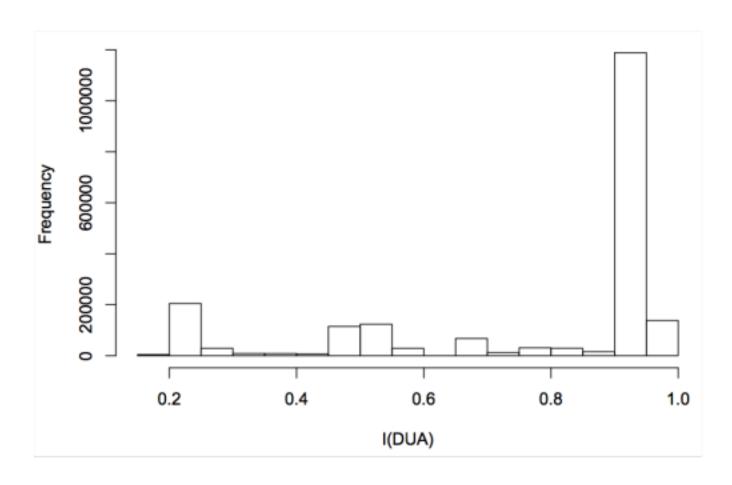


Fig. 8. Normalized DUA trace location

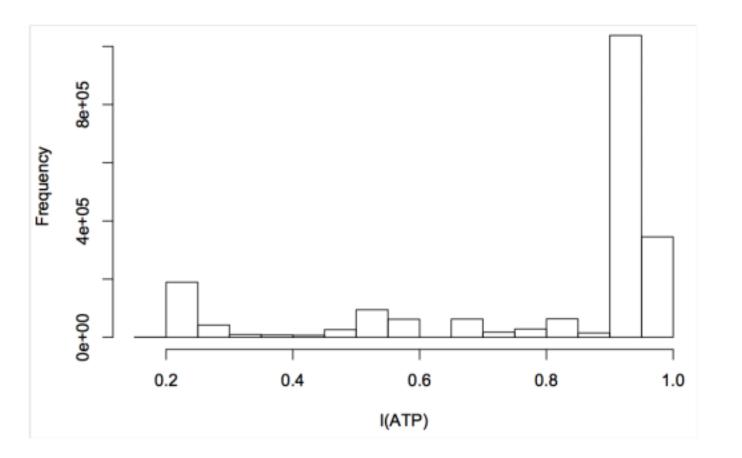


Fig. 9. Normalized ATP trace location



# Summary and future directions



### Summary

- Working system automates construction of large corpora for study and assessments
- Novel taint-based measures are key: liveness and TCN

### Future directions

- Continuous on-line competition to encourage self-eval
- Use in security competitions like Capture the Flag to reuse and construct challenges on-the-fly
- Assess and improve realism of LAVA bugs
- More types of vulnerabilities
- More interesting effects (exploitable ones)



